

ACFs are simple as much as infinite sets: An introduction to quantifier elimination

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- My goal is to convince that algebraically closed fields are simple as much as infinite sets (in the view of their syntax).
- I will introduce a notion of **quantifier elimination** (QE) from model theory.
- In the term of QE, the syntax of an algebraically closed field, as a field, is simple as like
 - an infinite set as a structure equipped with equality only,
 - a real closed field as an ordered field,
 - a differentially closed fields as a differential field.

- Fix a first order language \mathcal{L} , which is countable for simplicity. Let T be a complete \mathcal{L} -theory.
- Write x, y, z, \dots for tuples of variables.
- Let ω be the set of natural numbers.
- We say that a formula $\varphi(x)$ is **quantifier-free** if it has no quantifiers in φ .
- We say that T has **quantifier elimination** (QE) if for any formula $\varphi(x)$, there is a quantifier-free formula $\psi(x)$ such that

$$T \models \forall x(\varphi(x) \leftrightarrow \psi(x)),$$

that is, any formula is equivalent to a quantifier-free formula modulo T .

Example

Let T be the theory of real closed fields in the ordered ring language. Let

$$\varphi(a, b, c) \equiv a \neq 0 \wedge \exists x(ax^2 + bx + c = 0).$$

Then, φ is equivalent (modulo T) to the following quantifier-free formula

$$\psi(a, b, c) \equiv b^2 - 4ac \geq 0.$$

- Even though QE is defined syntactically, it has a semantic criterion, which is very useful.

Theorem

T has QE if and only if for a \aleph_0 -saturated and \aleph_1 -strongly homogeneous model \mathfrak{C} of T , the following holds: For an isomorphism $f : A \rightarrow B$ between finitely generated substructures A and B of \mathfrak{C} and $a \in \mathfrak{C}^1$, there is $b \in \mathfrak{C}^1$ such that the map $f \cup \{(a, b)\}$ is extended into an isomorphism between finitely generated substructures of \mathfrak{C} .

- A structure \mathfrak{C} is called \aleph_0 -saturated if the following holds: Let $\Sigma(x)$ be a countable set of $\mathcal{L}(\mathfrak{C})$ -formulae in the variable x of countable length. Suppose any finite subset Σ_0 of Σ has a solution in \mathfrak{C} . Then, there is a solution of Σ .
- A structure \mathfrak{C} is called \aleph_1 -strongly homogeneous if for any tuples \bar{a} and \bar{b} of elements in \mathfrak{C} of countable length,

$$\bar{a} \equiv \bar{b} \Rightarrow (\exists \sigma \in \text{Aut}(\mathfrak{C}))(\sigma(\bar{a}) = \bar{b}),$$

where $\bar{a} \equiv \bar{b}$ means $\mathfrak{C} \models \varphi(\bar{a}) \Leftrightarrow \mathfrak{C} \models \varphi(\bar{b})$ for all formulas $\varphi(x)$. 5 / 20

- Let $\mathcal{L} = \emptyset$ and let T be the theory of infinite sets.
- Then, T has QE.
- Let \mathcal{C} be a \aleph_0 -saturated and \aleph_1 -strongly homogeneous infinite set.
- Note that any subset of \mathcal{C} is a substructure because there are no function symbols.
- Let $f : A \rightarrow B$ be an isomorphism between finite subsets of \mathcal{C} with $|A| = |B| = n < \omega$.
- That is, f is just a bijection between A and B .
- Take $a \in \mathcal{C}$ arbitrary. If $a \in A$, then $f \cup \{(a, f(a))\}$ does work.
- Suppose $a \notin A$. Since \mathcal{C} is **infinite** and B is **finite**, there is $b \in \mathcal{C} \setminus B$. Then, the map $f \cup \{(a, b)\}$ does work.

- Let $\mathcal{L} = \{<\}$ and let T be the theory of linear orders without endpoints.
- Then, DLO has QE.
- Let \mathcal{C} be a \aleph_0 -saturated and \aleph_1 -strongly homogeneous DLO.
- Note that any subset of \mathcal{C} is a substructure because there are no function symbols.
- Let $f : A \rightarrow B$ be an isomorphism between finite subsets of \mathcal{C} with $|A| = |B| = n < \omega$.
- That is, f is an increasing bijection between A and B .
- Write $A := \{a_0 < a_1 < \dots < a_{n-1}\}$ and $B := \{b_0 < b_1 < \dots < b_{n-1}\}$ with $b_i = f(a_i)$.

- Take $a \in \mathcal{C} \setminus A$ arbitrary.
- Then, there are essentially $(n + 1)$ -many cases:
 - ① $a < a_0$.
 - ② For some $0 \leq i < n - 1$,

$$a_i < a < a_{i+1}.$$

- ③ $a > a_{n-1}$.
- Suppose $a_0 < a < a_1$.
 - Then, since \mathcal{C} is **dense**, there is b such that $b_0 < b < b_1$, and the map $f \cup \{(a, b)\}$ does work.
 - For the first and third cases, it comes from the fact that \mathcal{C} has **no endpoints**.

- Let $\mathcal{L}_{ring} = \{+, \cdot, 0, 1\}$ be the ring language and ACF_p be the theory of algebraically closed fields of characteristic p .
- Then, ACF_p has QE.
- Let \mathcal{C} be a \aleph_0 -saturated and \aleph_1 -strongly homogeneous model of ACF_p .
- For a subset A of \mathcal{C} , the substructure generated by A is the field generated by A .
- Let $f : A \rightarrow B$ be an isomorphism between finitely generated subfields of \mathcal{C} .
- The isomorphism f can be extended into an isomorphism between the algebraic closure of A and B .
- WLOG, we may assume that A and B are algebraically closed.

- Take $a \in \mathfrak{C} \setminus A$ arbitrary. Then a is transcendental over A .
- We can take $b \in \mathfrak{C}$ which is transcendental over B because \mathfrak{C} is \aleph_0 -saturated and B is countably generated.
- Then, there is an isomorphism

$$f' : A(a) \cong_A A(X) \cong B(X) \cong_B B(b), a \mapsto a$$

extending f .

- The similar process works for the theory DCF_0 of differentially closed fields of characteristic 0 in the differential ring language.
- In this case, we work with
 - Differential polynomials analogous to polynomials,
 - Differential ideals analogous to ideals,
 - The Kolchin topology, which is Noetherian, analogous to the Zariski topology.

- **QE** is very much dependent on **the choice of a language**.
- Consider the field \mathbb{R} of reals.
- Let T_1 be the theory of \mathbb{R} in the ring language $\mathcal{L}_1 := \mathcal{L}_{ring}$ and let T_2 be the theory of \mathbb{R} in the ordered ring language $\mathcal{L}_2 := \mathcal{L}_{ring} \cup \{<\}$.
- In \mathbb{R} , $<$ is definable in the ring language, that is,

$$\mathbb{R} \models \forall x, y (x < y \leftrightarrow \exists z (y = z^2 + x)).$$

- So, \mathbb{R} has the exactly same definable sets or the same ‘expressing’ power in both languages of \mathcal{L}_1 and \mathcal{L}_2 . More generally, the same thing holds for all **real closed** fields.
- A field F is called **real closed** if
 - it is formally real, that is, -1 is not a sum of squares,
 - any polynomial over F of odd degree has a zero in F .

- We will show that T_1 has no QE in \mathcal{L}_1 but T_2 has QE in \mathcal{L}_2 .
- Let \mathfrak{C} be a real closed field which is \aleph_0 -saturated and \aleph_1 -strongly homogeneous so that it is a model of T_1 and of T_2 .
- T_1 has no QE in $\mathcal{L}_1 = \mathcal{L}_{ring}$:
- Consider an ring isomorphism

$$f : \mathbb{Q}(\sqrt{2}) \rightarrow \mathbb{Q}(-\sqrt{2}), \sqrt{2} \mapsto -\sqrt{2}.$$

- Take $a = \sqrt[4]{2} \in \mathfrak{C}$. Then, we can not find $b \in \mathfrak{C}$ such that there is an ring isomorphism between finitely generated subfields of \mathfrak{C} , extending $f \cup \{(a, b)\}$.
- Why? Suppose there is such a 'b'.

$$a^2 = \sqrt{2} \Rightarrow b^2 = -\sqrt{2}.$$

- In the **real closed** field \mathfrak{C} ,

$$0 < b^2 = -\sqrt{2} < 0,$$

a contradiction.

- QE of ACF implies Chevalley's theorem on constructible sets.

Theorem

The set of *constructible* sets on \mathbb{C} is closed under taking projection.

- An algebraic subset of \mathbb{C}^n is a zero of polynomial equations over \mathbb{C} .
- A subset of \mathbb{C}^n is called *constructible* if it is a boolean combination of algebraic subsets.
- Chevalley's theorem says that given a constructible subset A of \mathbb{C}^{n+1} , the projection $\pi[A]$ is also constructible, where $\pi : (x_0, \dots, x_n) \mapsto (x_1, \dots, x_n)$.

- By definition, a subset of \mathbb{C}^n is constructible if and only if it is definable by a quantifier-free formula over \mathbb{C} .
- Let $A \subseteq \mathbb{C}^{n+1}$ be constructible.
- So, there is a quantifier-free formula $\varphi(x_0, \dots, x_n)$ such that

$$A = \{\bar{a} \in \mathbb{C}^{n+1} : \mathbb{C} \models \varphi(\bar{a})\}.$$

- Then,

$$\pi[A] := \{(b_1, \dots, b_n) \in \mathbb{C}^n : \exists x_0 \in \mathbb{C} ((x_0, b_1, \dots, b_n) \in A)\}.$$

- That is, for $\psi(x_1, \dots, x_n) \equiv \exists x_0 \varphi(x_0, x_1, \dots, x_n)$,

$$\pi[A] := \{\bar{b} \in \mathbb{C}^n : \mathbb{C} \models \psi(\bar{b})\}.$$

- By QE, ψ is equivalent to a quantifier-free formula, and so $\pi[A]$ is again constructible.

- Hilbert's 17th problem (theorem) says that given a polynomial $p(T) \in \mathbb{R}[T]$ with $|T| \geq 1$, if $p(a) \geq 0$ for all $a \in \mathbb{R}$, then p is a sum of squares of rational polynomials in $\mathbb{R}(T)$.
- It was first proved by Artin in 1927.

Example

Motzkin provided an example of polynomial having non-negative values for reals but not sum of squares of polynomials over \mathbb{R} :

$$x^4y^2 + x^2y^4 + 1 - 3x^2y^2 = \frac{x^2y^2(x^2 + y^2 + 1)(x^2 + y^2 - 2)^2 + (x^2 - y^2)^2}{(x^2 + y^2)^2}.$$

- Using QE of RCF , we will give a model theoretic proof of Hilbert's 17th problem (by Robinson in 1955).
- QE of RCF implies that RCF is **model-complete**:
- For $M, N \models RCF$ with $M \subseteq N$, then M is an elementary substructure of N , denoted by $M \prec N$, that is, for any formula $\varphi(x)$ and $a \in M^{|x|}$,

$$M \models \varphi(a) \Leftrightarrow N \models \varphi(a).$$

- Suppose there is a polynomial $p(T) \in \mathbb{R}[T]$ with $T = (T_0, \dots, T_{n-1})$ such that
 - $p(a) \geq 0$ for all $a \in \mathbb{R}^n$,
 - $p \neq q_0^2 + \dots + q_m^2$ for all $q_0, \dots, q_m \in \mathbb{R}(T)$.

Fact

For a field F and $a \in F$, suppose -1 is not a sum of squares in F and a is not a sum of squares in F . Then, there is a linear order $<$ on F such that $(F, <)$ is an ordered field with $a < 0$.

- By the above fact, there is a linear order $<'$ on $\mathbb{R}(T)$ such that $(\mathbb{R}(T), <')$ is an ordered field with $p(T) <' 0$.
- Note that $<'$ is extending the linear order $<$ on \mathbb{R} because for any real number a , either a or $-a$ is a square of real number.

Fact

Any formally real field F has a real closure (**unique** up to isomorphism over F), which is a real closed algebraic extension of F .

- Let $(F, <')$ be the real closure of $(\mathbb{R}(T), <')$, extending $(\mathbb{R}, <)$
- By model-completeness, $(F, <')$ is an elementary extension of $(\mathbb{R}, <)$.
- By the choice of $p \in \mathbb{R}[T]$, we have that

$$(\mathbb{R}, <) \models \forall x (p(x) \geq 0).$$

- Since $\mathbb{R} \prec F$, we have that

$$(F, <') \models \forall x (p(x) \geq 0),$$

- Since $T_0, \dots, T_{n-1} \in \mathbb{R}(T) \subseteq F$, for $T := (t_0, \dots, t_{n-1})$

$$(F, <') \models 0 \leq' p(T),$$

which contradicts with $p(T) <' 0$.

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Thank you for your listening
Happy Logic Day